Computational Geometry

Wrap-Up

Thomas Bläsius

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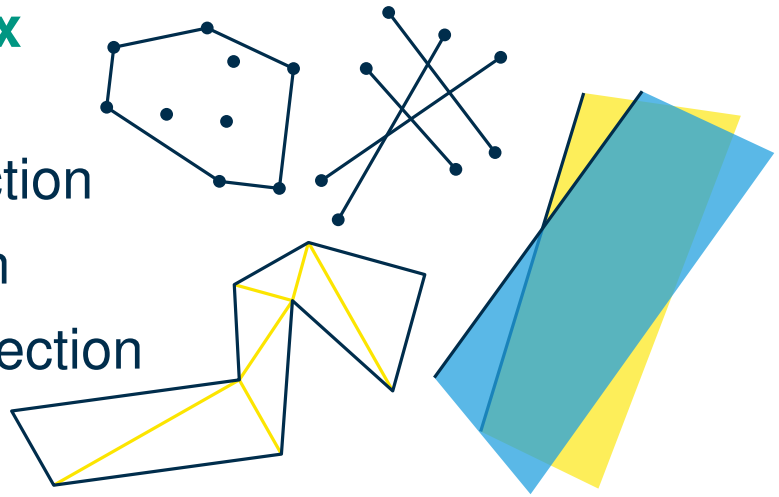
- the questions on the following slides may serve as a starting point for this

Overview

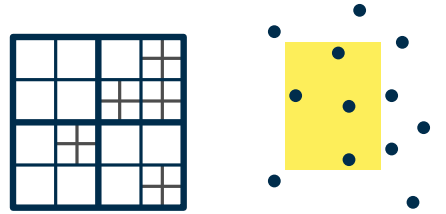
Basic Toolbox

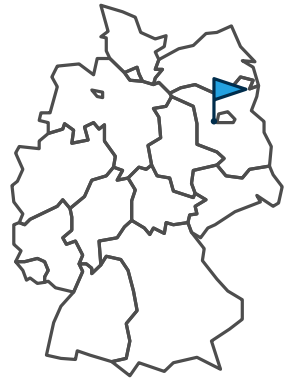
- convex hull
- line intersection
- triangulation
- plane intersection



Geometric Data Structures

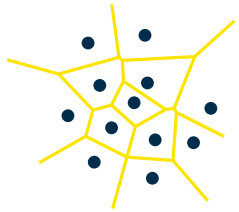
- orthogonal range searching
- space partitioning
- point location

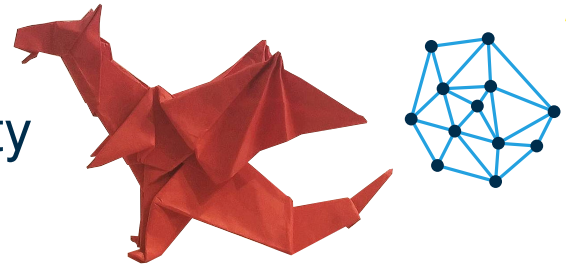




Advanced Toolbox

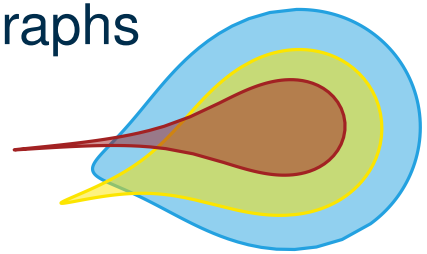
- Voronoi diagrams
- Delaunay triangulations
- origami
- complexity

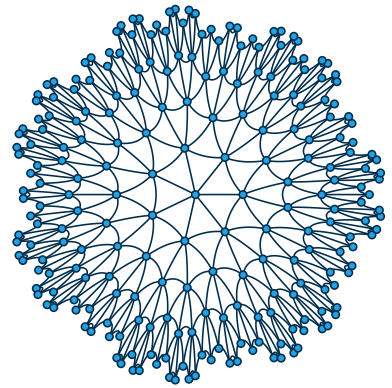




Related Topics

- What is geometry?
- hyperbolic geometry
- geometric graphs

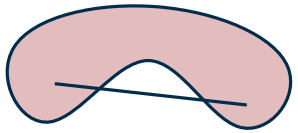


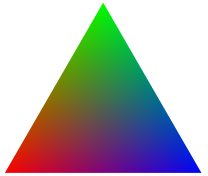


Convex Hull

Problem Definition

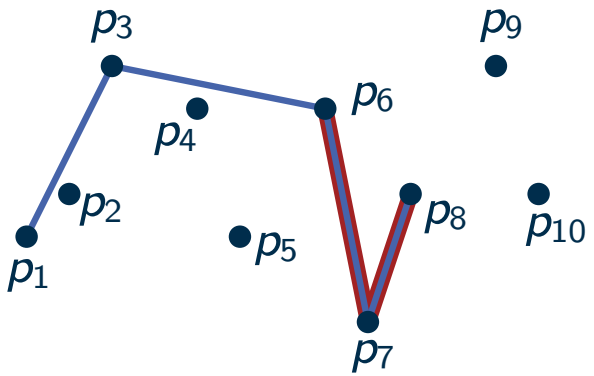
- What is a convex set? Different definitions?
- What problem did we consider?

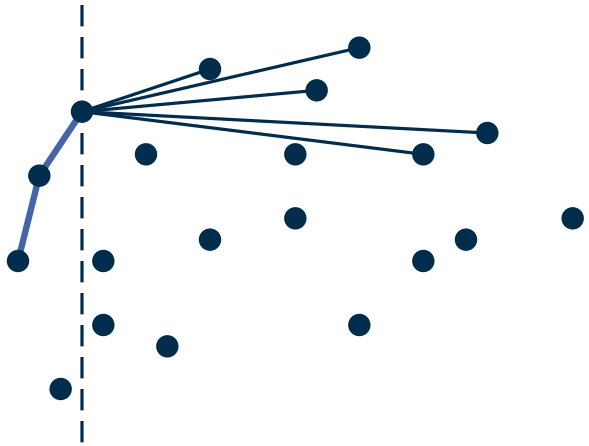




Algorithms: Andrews Algorithm (Graham Scan)/Gift Wrapping

- How does the algorithm work?
- Why is it correct?
- What is the running time? Why?
- Is it robust?





Lower Bounds

- Can it be done faster? Why not?

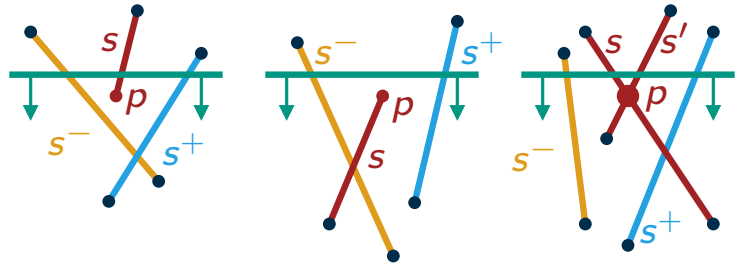
~~Convex Hull On A Simple Polygon~~

- ~~Can this be done faster?~~
- ~~Why? How do you use the given polygon?~~

Segment Intersection & Sweep-Line

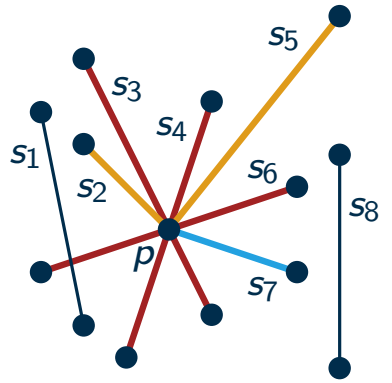
Intersection Of Line Segments

- What is the general approach?
- What special cases do we need to consider?
- How do we handle multiple events at the same point?



Running Time

- What is the running time? How can we show that?



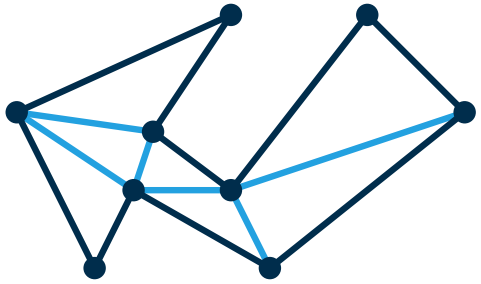
Memory Consumption

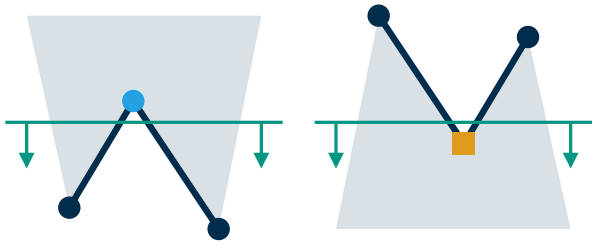
- Is linear memory sufficient? How do we achieve that?
- Do we need $\Omega(n^2)$ memory, if we are not careful?

Doubly-Connected Edge List

- What is it? What is it good for?
- What is the connection to segment intersection?

Triangulating Polygons



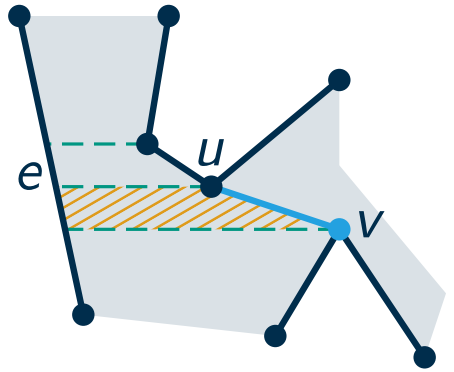


Basics

- What is the idea?
- What does y-monotone mean?
- What are split and merge vertices and how do they relate to y-monotonicity?

Eliminating Split Vertices

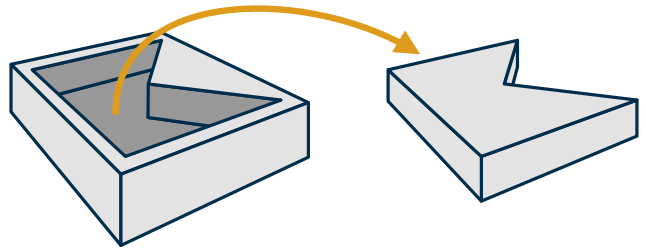
- How can we get rid of split vertices?
- What is the role of the helper?
- Why are the inserted edges crossing-free?
- How does this yield an algorithm? What events are there? How do we handle them?
- What running time do we get?



Triangulating

- What is the idea for triangulating y-monotone polygons?

Linear Programs

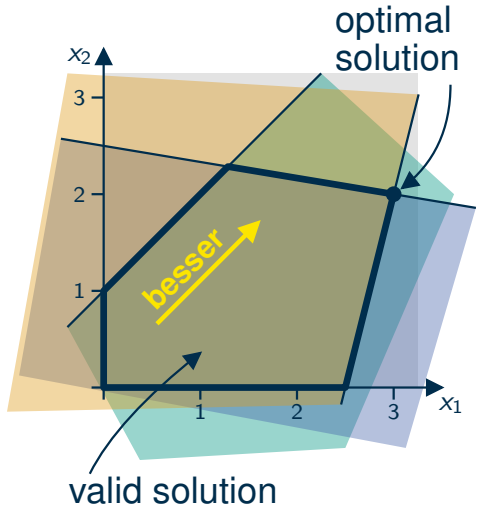


Molds

- What is the mold problem? How can it be solved using a 2D-LP?

Linear Programming

- What is an LP? What is the connection to geometry?

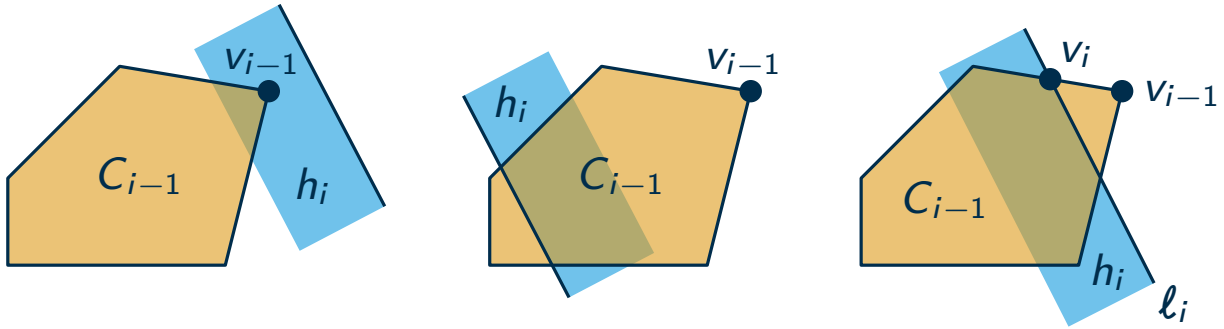


Algorithm 1: Divide And conquer

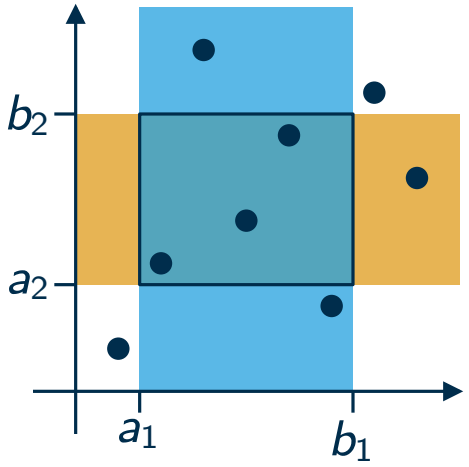
- What is the idea? What running time do we get?

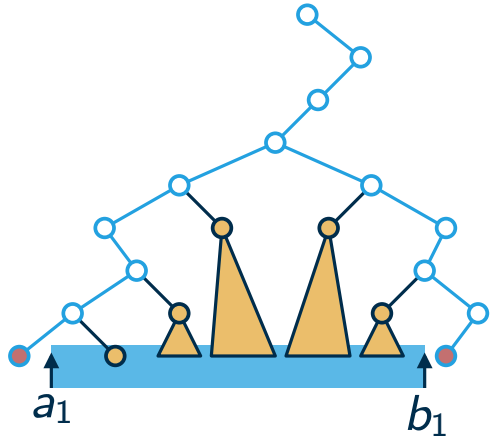
Algorithm 2: Incremental Approach

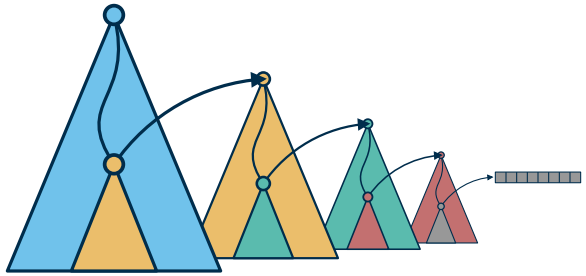
- What is the idea?
- How does it reduce to a 1D-LP in each step?
- Why does a random order help?
- What is the running time? Why?
- How does the backwards-analysis help?



Orthogonal Range Queries







Basics

- What is the setting?

Range Trees

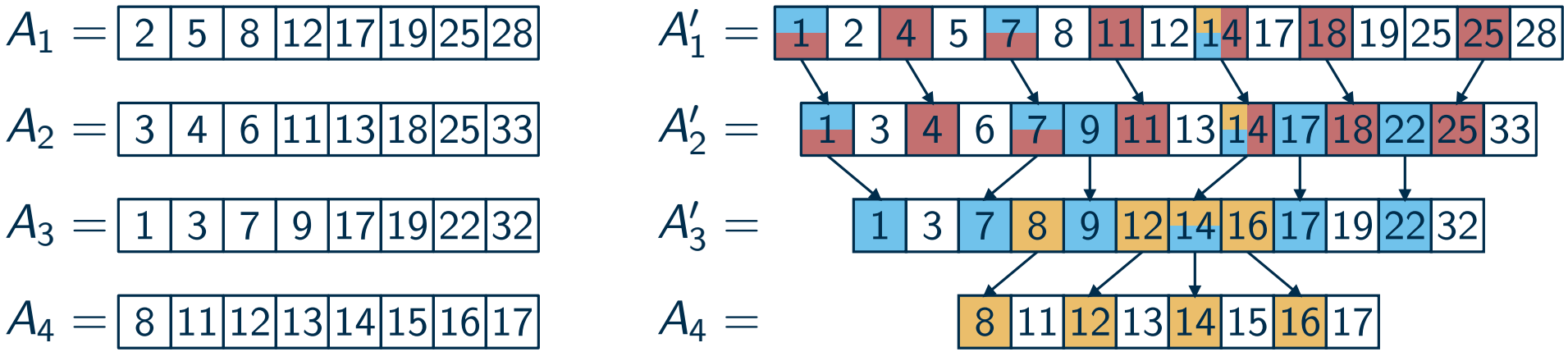
- What is a 2D-range tree?
- How can we answer queries and how long does it take?
- How can we build it and how long does it take? How much memory?
- How does it generalize to more dimensions?

Improved Queries

- How can we save $\log n$ searches?
- How does this generalize to fractional cascading?
- How to answer queries of the form $[-\infty, b_2] \times [-\infty, b_3]$? And $[a_1, b_1] \times [-\infty, b_2] \times [-\infty, b_3]$?
- Can we get rid of the ∞? What is the impact on running time and memory?

Interlude: Fractional Cascading

Core Insight: Searching for the same x in ℓ sorted arrays of size n can be done in $\ell + \log n$ time.
(linear preprocessing and memory)



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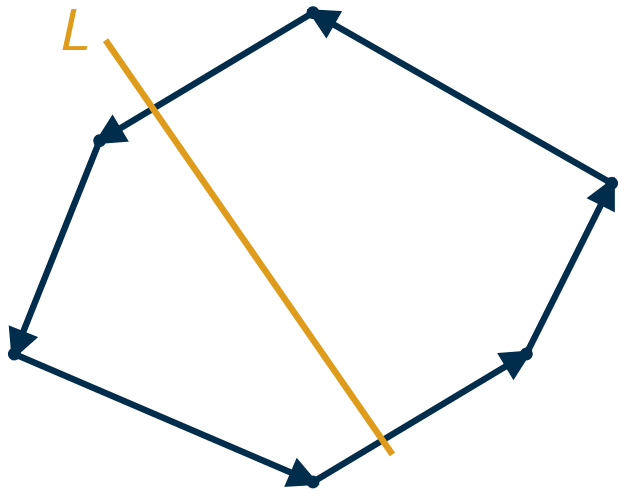

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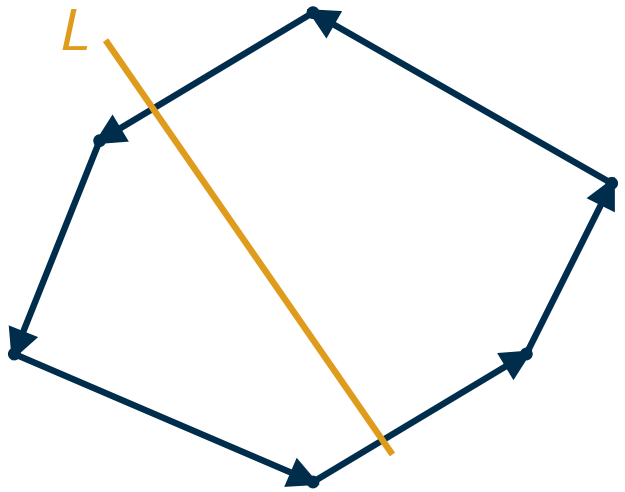

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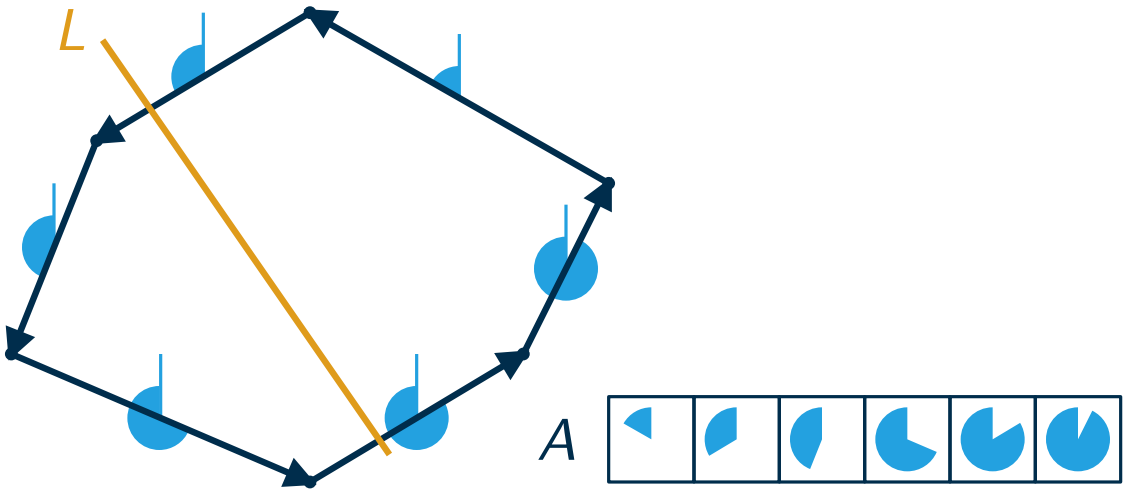

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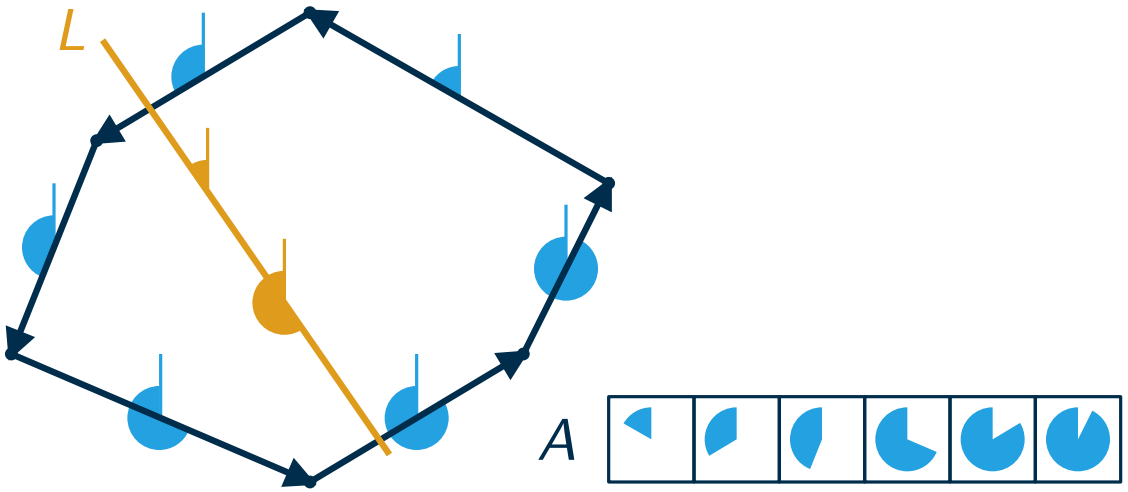

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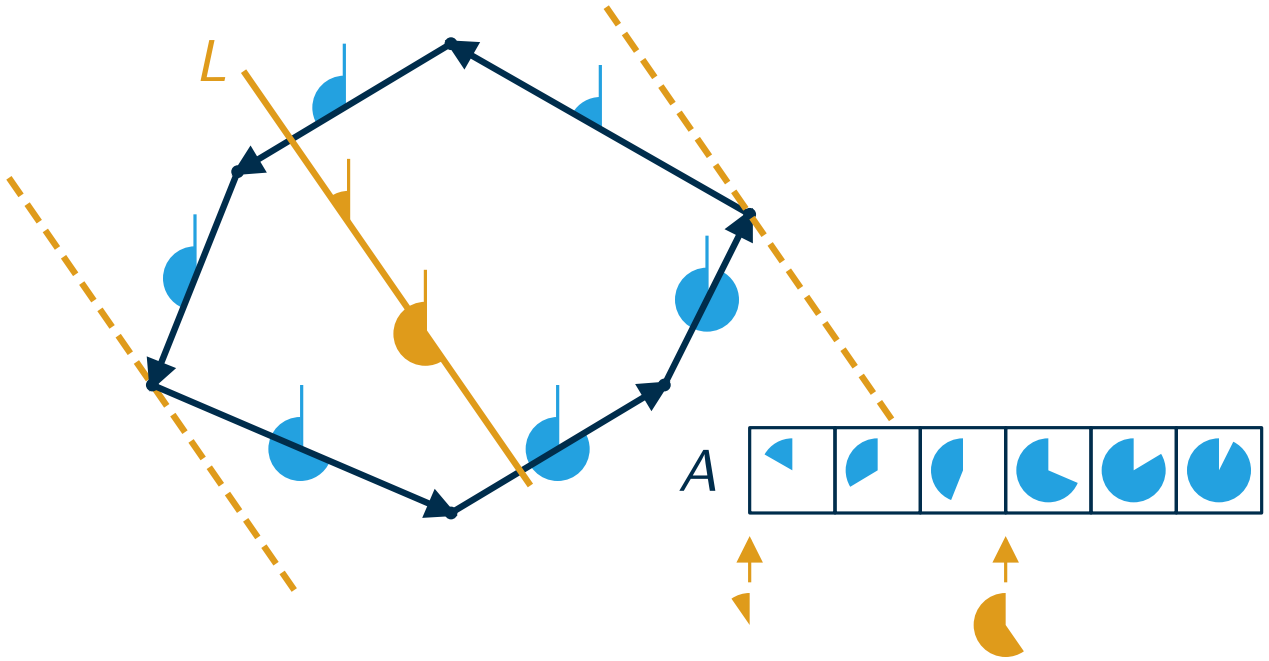

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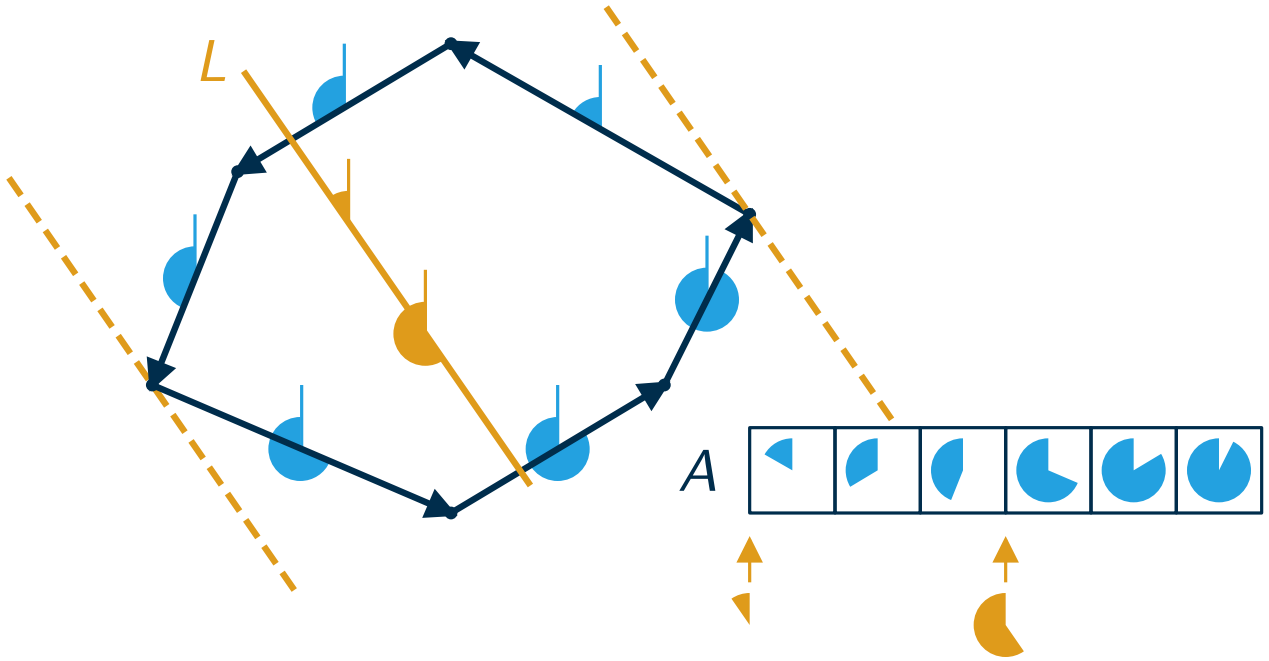

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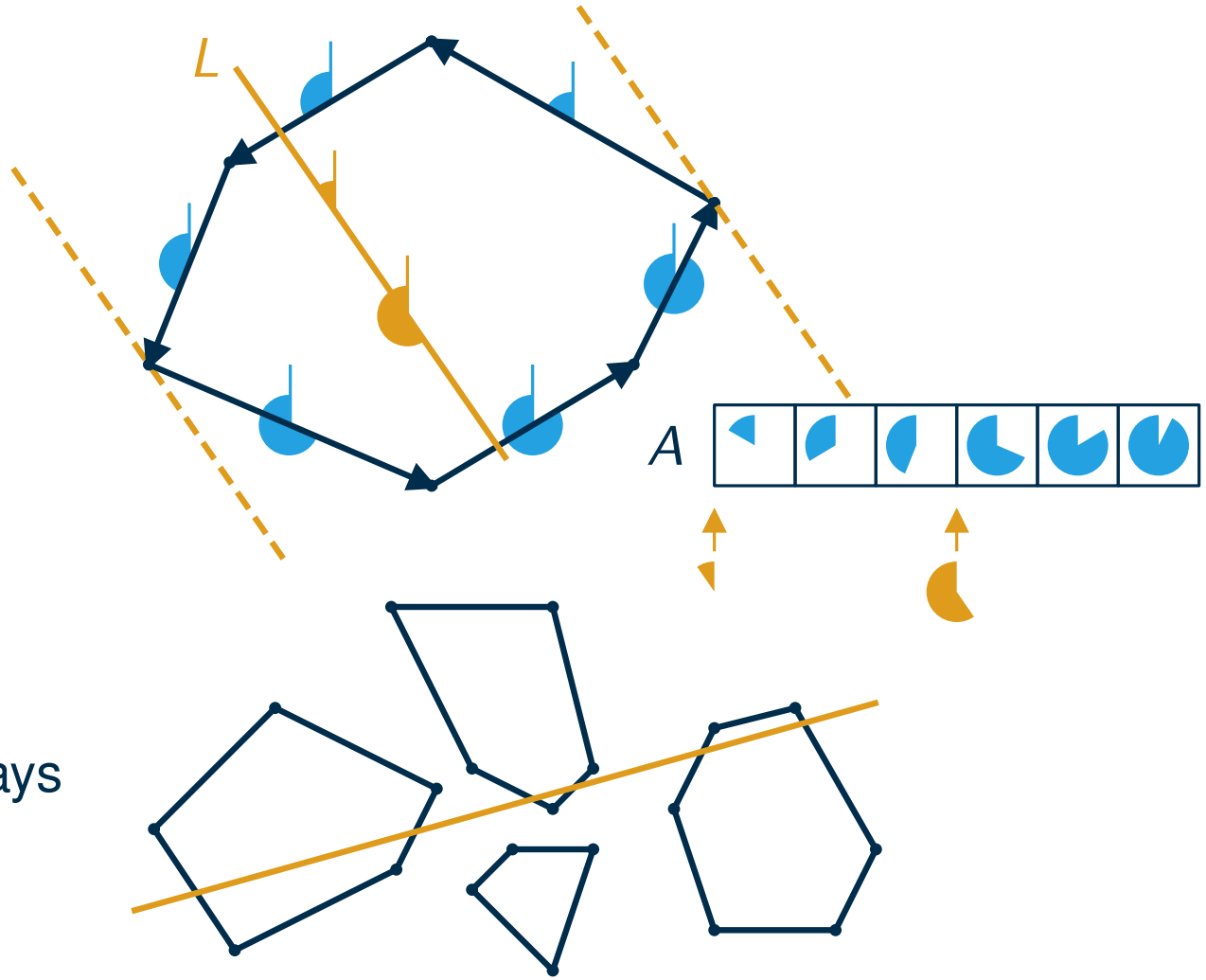

Now With Multiple Convex Polygons

- ℓ convex polygons $\rightarrow$ search in ℓ angle arrays

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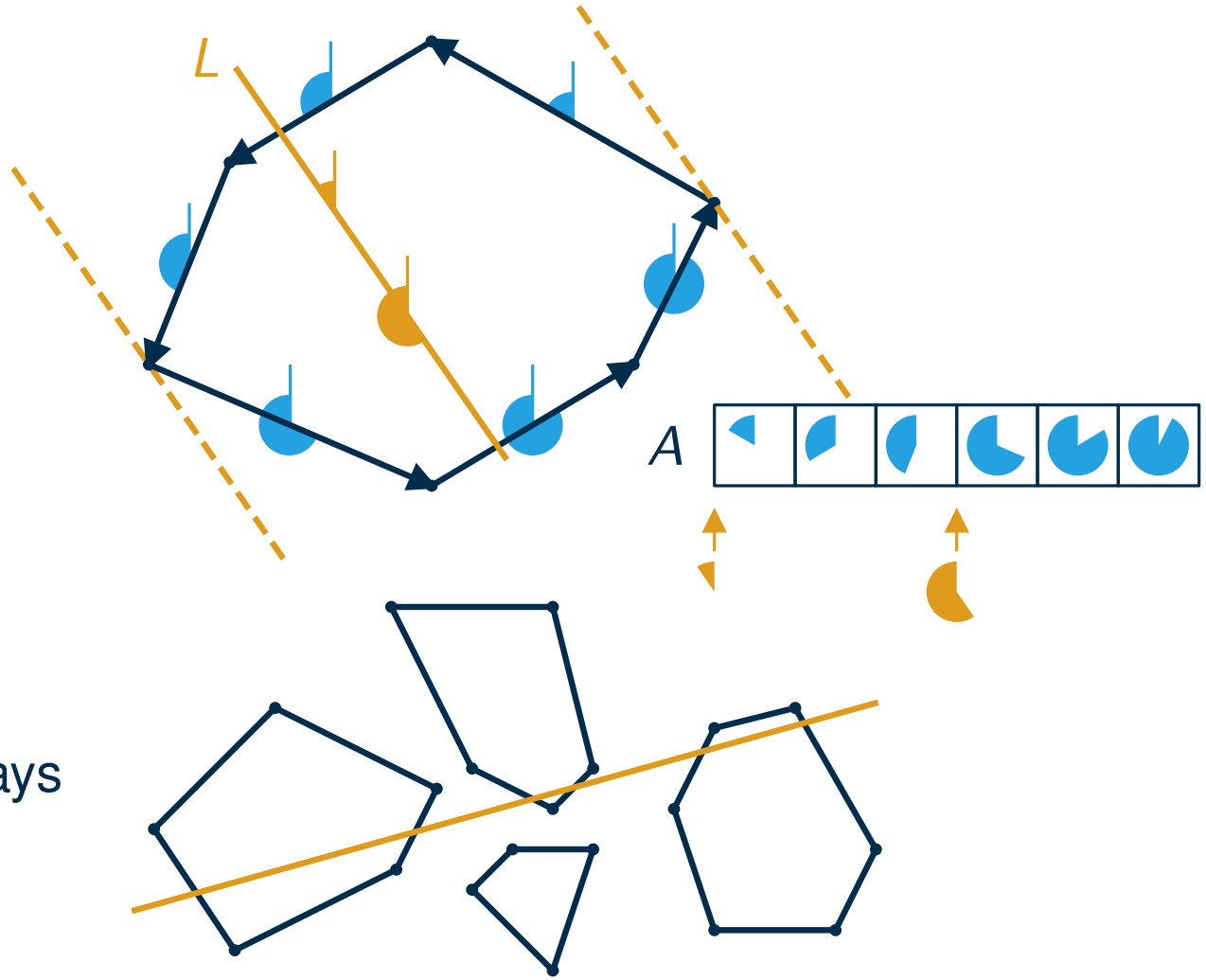

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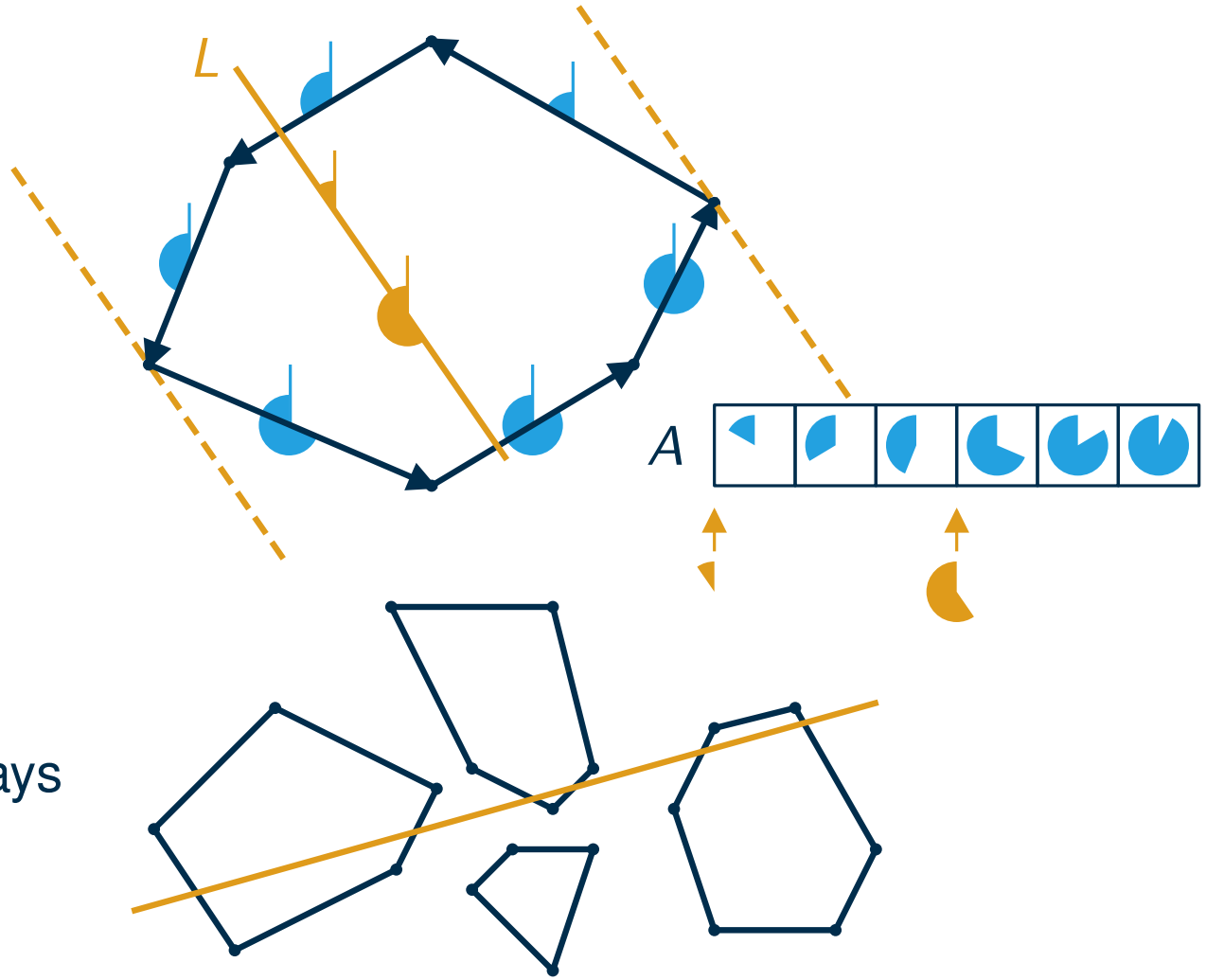

Now With Multiple Convex Polygons

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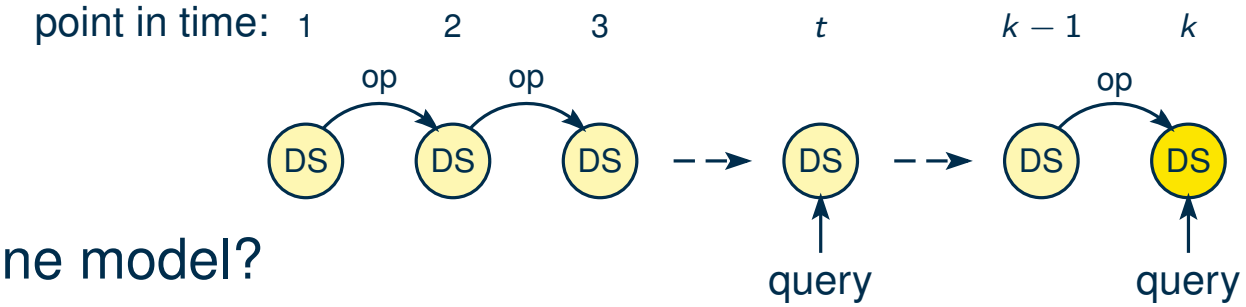
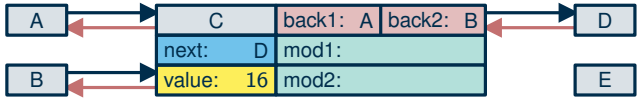
- fractional cascading $\rightarrow O(\ell + \log n)$
- for nested polygons: pruning possible

Time Travel and Point Location

Persistence

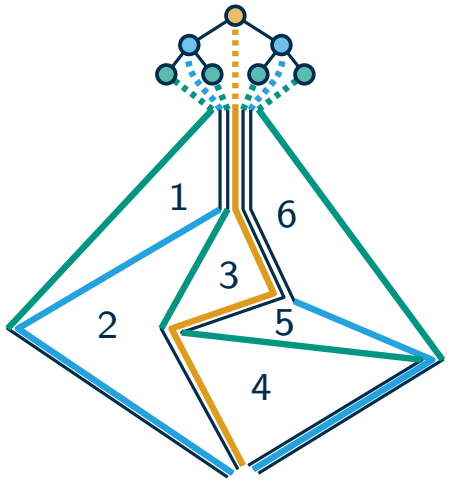


- What is a persistent data structure?
- What types of persistence are there?
- What is a data structure in the pointer machine model?
- How do we make such a data structure persistent?
  - What are atomic operations? How do we make them persistent?
  - Why are cascading recursive calls not too expensive?
  - Why do we need bounded in-degree?



Point-Location

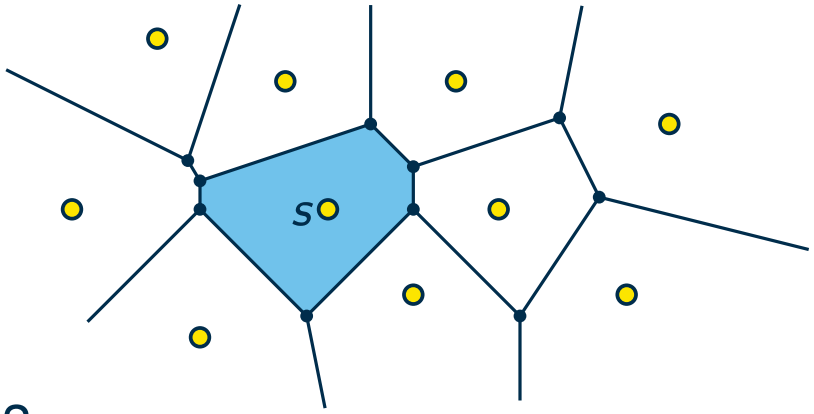
- What is the Problem? How can we solve it using persistence?
- What running time and memory consumption do we get?
- What is an alternative solution? How does fractional cascading help?



Voronoi Diagrams

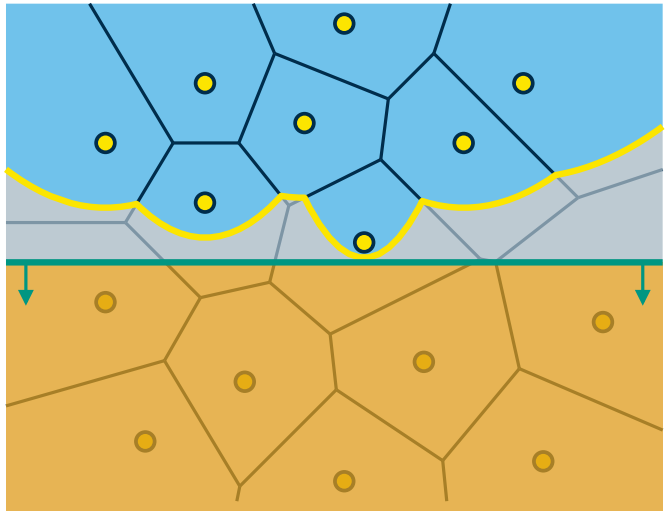
Basics

- What is the Voronoi diagram?
- Which points are vertices of the Voronoi diagram?



Sweep-Line Algo

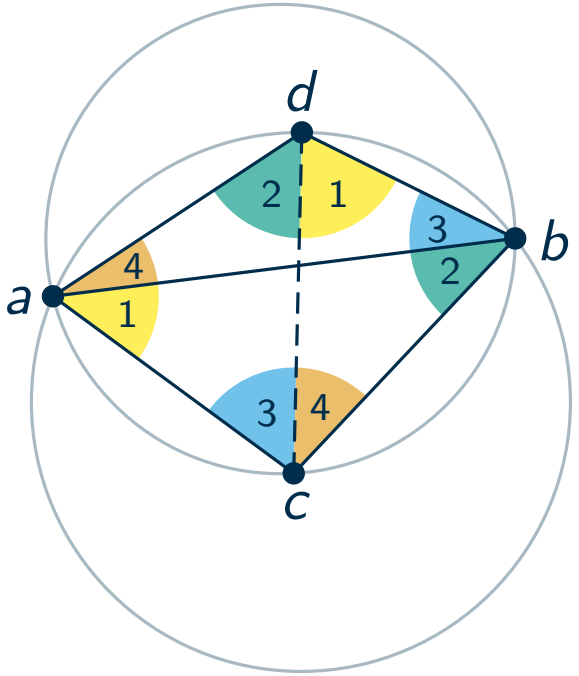
- What is different compared to previous sweep-line algorithms?
- What is the beach line? What is its shape and why?
- In what sense does the beach line only change at discrete places?
- What are these places? What are our events?
- How do we handle these events?
- How do we get the Voronoi diagram in the end?
- What about special cases?



Delaunay Triangulation

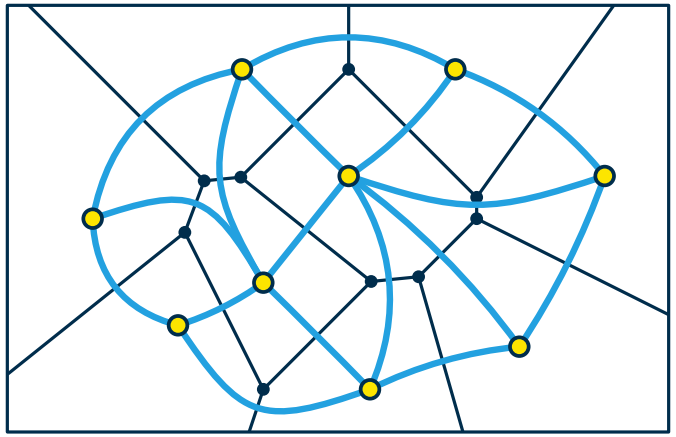
Improving Triangulations

- What is the angle vector of a triangulation?
- How can we iteratively improve a given triangulation?
- What are forbidden edges?
- Which edges are forbidden and what is the connection to circles?
- How does (the generalization of) Thales' theorem help?

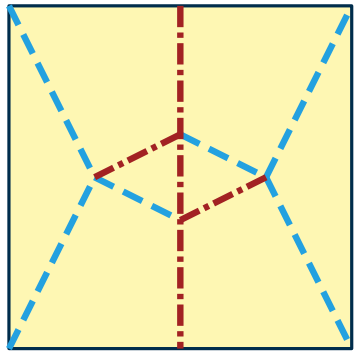


(Locally) Optimal Triangulations

- What is the Delaunay triangulation? Is it unique?
- What is the relation to the Voronoi diagram?
- Why is the relation the way it is?



Origami

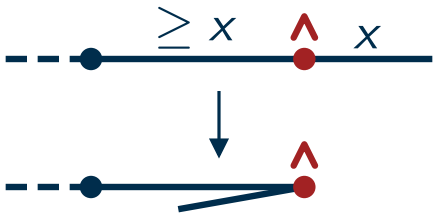


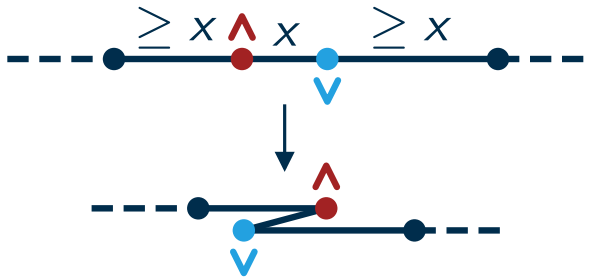
Flat Foldability

- What is a crease pattern? What is a mountain/valley pattern?

1D-Case

- What are crimp and end-fold?
- Are these reduction rules safe? Why?
- Are they always applicable? Why?
- How does this yield an algorithm?



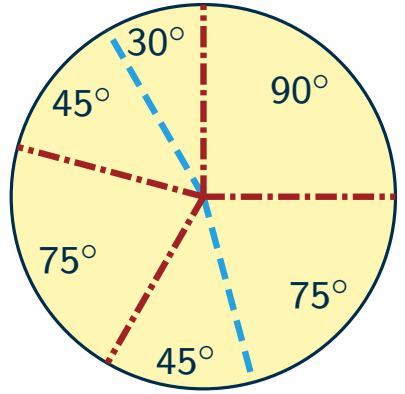


2D-Case: One Vertex

- What is the relation to the 1D-case?
- Which crease and mountain/valley patterns are flat foldable?

Multiple Vertices

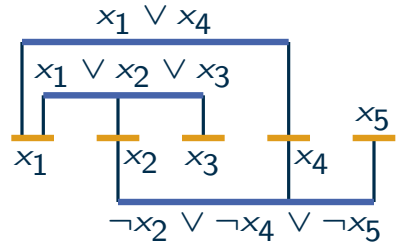
- What is local flat foldability?
- What is the idea for fold-and-cut?



Hard Problems

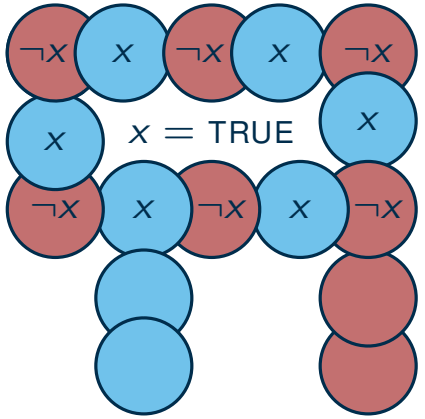
Reductions From Hard SAT-Variants

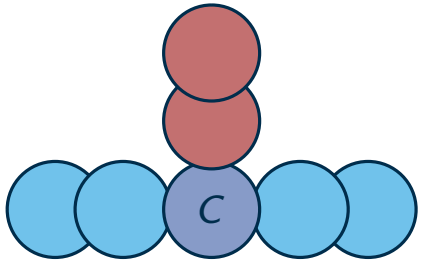
- Which SAT-variants are useful for geometric problems?
- What do we need to do in a reduction? What do we then need to show?



Example 1: Proportional Symbol Maps

- What is the problem? What variants exist?
- How do we prove hardness?

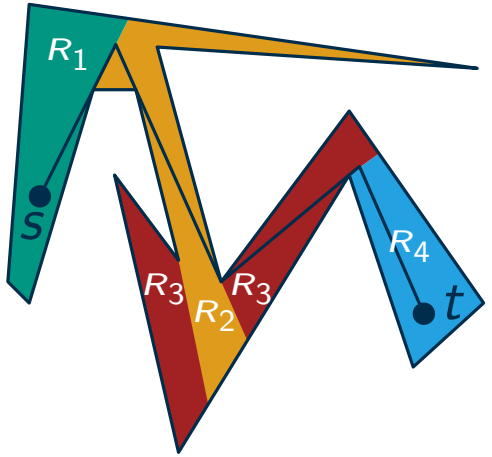




Example 2: Unit Disk Graph Recognition

- What is the problem?
- How do we prove hardness?
- Why do we need a splitter gadget?
- Why can the graph not be drawn completely differently?
- What is the complexity class $\exists\mathbb{R}$?

Searching In Sublogarithmic Time

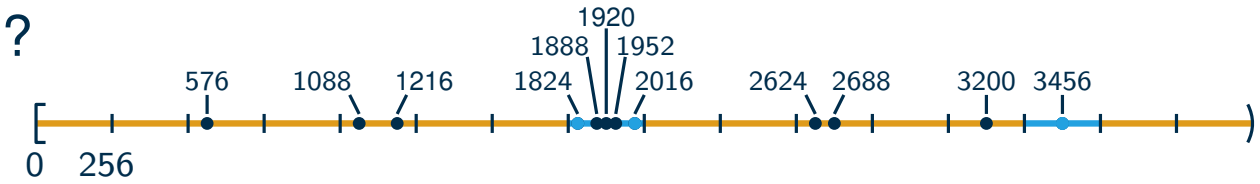


Model Of Computation

- What are RAM, real RAM, and word RAM?
- Which operations are supported in constant time?

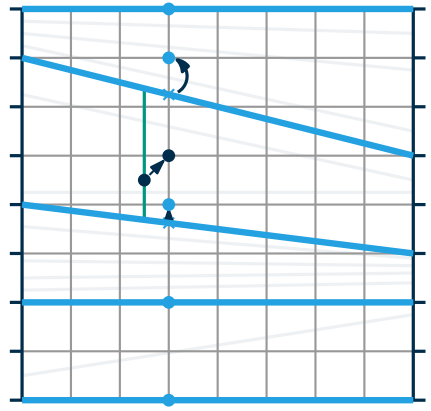
Real RAM

- Why is it useful for us? What are the pitfalls?



Word RAM: Searching

- Which two ideas do we use to speed-up the search?
- How do we branch in every search step? How high is the decision tree?
- How do we make the decision in which child do descend? In $O(1)$ time?



Word RAM: 2D-Search

- How does it relate to the 1D-case? What are the difficulties?
- What is the core idea to solve them?

Hyperbolic Geometry

Axiomatic Approach

- What is it?
- How to define geometric things without intuition?
- How are the basic terms filled with life?
- What is the Euclidean / absolute / hyperbolic plane?

Hyperbolic Plane

- What is it? Which statements still hold?
- How can I get some intuition? What is the Poincaré disk?

Hyperbolic Uniform Disk Graphs

- What is different compared to Euclidean UDGs?
- What properties do we have in the strongly hyperbolic setting?
- What can we do with it?

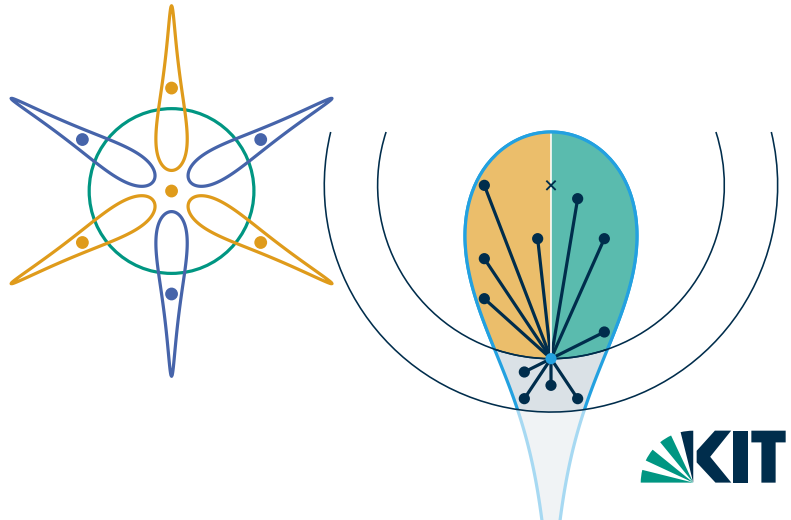
Axiom Group I: Incidence
two points define a line; every line contains two points; there are three non-collinear points

Axiom Group II: Distance
distance is a metric; tightness of triangle inequality if and only if collinear

Axiom Group III: Order
there is a point in every direction with every distance; lines split the plane into half planes

Axiom Group IV: Motion
two motions that map segments of equal length onto each other (preserving orientation)

Axiom Group V: Euclidean Parallel Axiom
line ℓ and point $P \notin \ell \Rightarrow$ at most one line through P parallel to ℓ



Good Luck!