

## **Probability & Computing**

#### Concentration





## **Expectation Management**

### What does it mean?

- "QuickSort has an *expected* running time of  $O(n \log(n))$ ."
- "The vertex has an *expected* degree of *c*."
- "In expectation there is one hair in my soup."

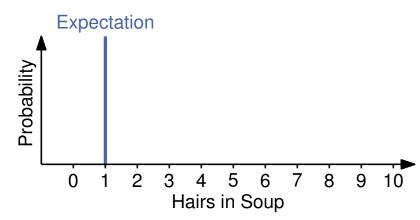
## Expectation

- The average of infinitly many trials
- How useful is that information in practice?
- Does not tell us much about the shape of the distribution
- Does not come with a level of certainty

## Concentration

- In practice, expectation is often a good start
- But for meaningful statements, we need to know how likely we are close to the exepcation

**Definition**: A **concentration inequality** bounds the probability of a random variable to deviate from a given value (typically its expectation) by a certain amount.



Knowing that the expected value is 1 hair:

How likely is it that I get at least 10? Not at all Somewhat

How likely is it that I get less than 2? Extremely Somewhat

## **Markov's Inequality**



#### **About Markov**

- Andrei "The Furious" Andreyevich Markov (Russian mathematician)
- Unhappy with the state of living at the time (1921)
- Informed St. Petersburg Academy of Sciences that he could not attend meetings due to not having shoes
- After getting shoes from the Communist Party he replied:

Finally, I received footwear. However, it is stupidly stitched together and does not accord with my measurements. Thus I cannot attend the meetings. I propose placing the footwear in a museum, as an example of the material culture of the current time.

"Shape, The hidden geometry of absolutely everything", Jordan Ellenberg

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**Corollary**: Let X be a non-negative rand. var. and a > 0. Then,  $\Pr[X \ge a \cdot \mathbb{E}[X]] \le 1/a$ .

• "In expectation there is one hair in my soup."

- How likely is it that I get at least 10?
- How likely is that I get less than 2?

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## Markov's Inequality

**Theorem (Markov's inequality)**: Let X be a non-negative random variable and let a > 0. Then,  $\Pr[X \ge a] \le \mathbb{E}[X]/a$ .

# Visual Proof $\mathbb{E}[X] = \mathbb{E}[X \mid X < a] \cdot \Pr[X < a] + \mathbb{E}[X \mid X < a] \cdot \Pr[X < a] + \mathbb{E}[X \mid X > a] \cdot \Pr[X > a] \cdot \Pr[X > a] \cdot \Pr[X > a]$

$$\begin{array}{l} \Pr[X \geq 10] \leq 1/10 \\ \Pr[X < 2] = 1 - \Pr[X \geq 2] \geq 1 - 1/2 = 1/2 \\ Oh \ no... \end{array}$$



## **Application: Unfair Coins**

The sum of 20 unfair  $\{0, 1\}$ -coin tosses:  $X \sim Bin(20, \frac{1}{5})$ What is the probability of getting at least 16 ones?

$$\Pr[X \ge 16] \le \mathbb{E}[X] / 16 = 0.25$$

How tight is that bound? Not very?

$$\Pr[X \ge 16] = \sum_{k=16}^{20} {20 \choose k} (\frac{1}{5})^k \cdot (1 - \frac{1}{5})^{20-k} \approx 0.0000000138$$

- Fair Coin
- A single  $\{0, 1\}$ -coin toss:  $Y \sim \text{Ber}(\frac{1}{2})$
- What is the probability of getting at least 1?
  - Clearly:  $\Pr[Y \ge 1] = \Pr[Y = 1] = \frac{1}{2}$
  - Markov:  $\Pr[Y \ge 1] \le \mathbb{E}[Y]/1 = \mathbb{E}[Y] = \frac{1}{2}$
- $\Rightarrow$  There is no better bound (that relies only on the expected value)



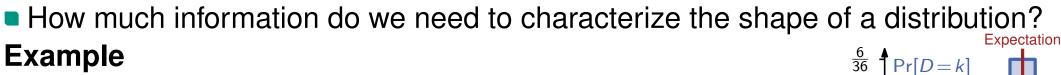
**Markov**: X non-negative, a > 0: Pr[ $X \ge a$ ]  $\le \mathbb{E}[X]/a$ .

Maybe it is just a weak bound?

There exists a random variable and an a > 0such that Markov's inequality is *exact*.

We need more information about the shape of the distribution!

## Characterizing the Shape of a Distribution

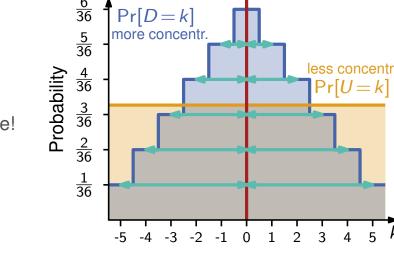


- X, Y independent fair die-rolls, D = X Y
- U uniform distribution over  $\{-5, -4, ..., 5\}$
- Consider all probabilities individually Tedious... We need to aggregate! **Expectation?** f(k) = k

 $\mathbb{E}[D] = \sum_{k} \Pr[D = k] \cdot k = 0$ Same value, different shapes  $\mathbb{E}[U] = \sum_{k} \Pr[U = k] \cdot k = 0 \text{ (also just seen with Markov: } \mathbb{E} \text{ not enough)}$ 

Problem: 
$$+ \& -$$
 terms cancel  
 $\Rightarrow$  Fix: absolute value  $f(k) = |k|$   
 $\mathbb{E}[|D|] = \sum_{k} \Pr[D = k] \cdot |k| \approx 1.945$   
 $\mathbb{E}[|U|] = \sum_{k} \Pr[U = k] \cdot |k| \approx 2.727$   
 $\mathbb{E}[D^2] = \sum_{k} \Pr[D = k] \cdot k^2 \approx 5.833$   
 $\mathbb{E}[|U|] = \sum_{k} \Pr[U = k] \cdot |k| \approx 2.727$   
 $\mathbb{E}[U^2] = \sum_{k} \Pr[U = k] \cdot k^2 = 10.0$   
Distance to  $\mathbb{E}$   
Squared distance to  $\mathbb{E}$ 

These are just expectations of functions of random variables!



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## Do you have a Moment?

#### **Expectation and Functions**

Random variable X taking values in a set S

• A function f, e.g. 
$$f(X) = X^1$$
,  $f(X) = |X|$ ,  $f(X) = X^2$ ,  $f(X) = \sqrt{X}$ ,  $f(X) = X^3$ ,  $f(X) = e^X$ 

 $\blacksquare \mathbb{E}[f(X)] = \sum_{x \in S} \Pr[X = x] \cdot f(x)$ 

These turn out to be particularly useful!

#### Moments

**Definition**: For random variable X and  $n \in \mathbb{N}$  the *n*-th raw moment is  $\mathbb{E}[X^n]$ .

• Just seen: For  $\mathbb{E}[X] = 0$ , this captures distances to  $\mathbb{E}[X]$  What if  $\mathbb{E}[X] \neq 0$ ?

**Definition**: For random variable X and  $n \in \mathbb{N}$  the *n*-th central moment is  $\mathbb{E}[(X - \mathbb{E}[X])^n]$ .

Just seen: the 2nd central moment captures squared distances to the expected value

 $\mathbb{E}[(X - \mathbb{E}[X])^2] = Var[X]$ 

The smaller the variance, the more concentrated the random variable

... and with Markov's help, we can turn that insight into a concentration inequality!



## Chebychev's Inequality

Markov's teacher! (Markov's inequality actually appeared earlier in Chebychev's works)

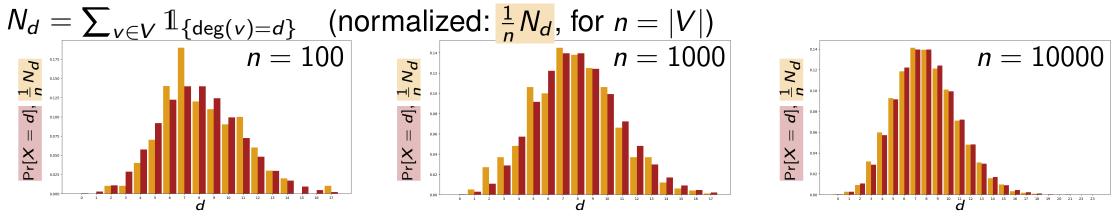
**Theorem (Chebychev's inequality)**: Let X be a random variable with finite variance and let b > 0. Then,  $\Pr[|X - \mathbb{E}[X]| \ge b] \le \operatorname{Var}[X]/b^2$ .

## **Application: ER – Degree Distribution**



#### Recap

- G(n, p): Start with *n* nodes, connect any two with fixed probability *p*, independently
- Probability distribution of the degree of a single node v:  $deg(v) \sim Bin(n-1, p)$
- For p = c/n with  $c \in \Theta(1)$  the degree of a vertex is approximately Poisson-distributed
  - Total variation distance of *X*, *Y* taking values in a set *S*:  $d_{TV}(X,Y) = \frac{1}{2} \sum_{x \in S} |\Pr[X = x] - \Pr[Y = x]|$
  - For  $\lambda = -n \log(1-p) = c + O(1/n)$  and  $X \sim \text{Pois}(\lambda)$  we have  $d_{TV}(\deg(v), X) = o(1)$
- Empirical distribution of the degrees of *all* vertices in a graph G = (V, E)



## **Application: ER – Degree Distribution**

**Theorem:** Consider a G(n, p) with p = c/n for constant c > 0. For  $\lambda = -n \log(1 - p)$ , let  $X \sim \text{Pois}(\lambda)$ . Then for all d > 0 and every  $\varepsilon > 0$  we have  $\lambda = c + O(1/n) \rightarrow c$  for  $n \rightarrow \infty$  $\lim_{n \rightarrow \infty} \Pr\left[\left|\Pr[X = d] - \frac{1}{n}N_d\right| \ge \varepsilon\right] = 0.$ 

#### Proof

Step 1: 
$$\Pr[X = d]$$
 is close to the expectation of  $\frac{1}{n}N_d$   $\lim_{n \to \infty} \left|\Pr[X = d] - \mathbb{E}\left[\frac{1}{n}N_d\right]\right| = 0 \checkmark$ 

$$\left|\Pr[X = d] - \mathbb{E}\left[\frac{1}{n}N_d\right]\right| = \left|\Pr[X = d] - \Pr[\deg(v) = d]\right| \le \sum_{d \ge 0} \left|\Pr[X = d] - \Pr[\deg(v) = d]\right|$$

$$= \frac{1}{n}\mathbb{E}[N_d]$$

$$= \frac{1}{n}\mathbb{E}[\sum_{v \in V} \mathbb{1}_{\{\deg(v) = d\}}]$$

$$= \frac{1}{n}\sum_{v \in V} \mathbb{E}[\mathbb{1}_{\{\deg(v) = d\}}]$$

$$= \frac{1}{n}\sum_{v \in V} \Pr[\deg(v) = d]$$

$$= \Pr[\deg(v) = d]$$
Step 2:  $\frac{1}{n}N_d$  is concentrated
$$\lim_{n \to \infty} \Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_d\right] - \frac{1}{n}N_d\right| \ge \varepsilon\right] = 0$$



$$\Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_{d}\right]-\frac{1}{n}N_{d}\right|\geq\varepsilon\right]\leq \operatorname{Var}\left[\frac{1}{n}N_{d}\right]/\varepsilon^{2}$$

$$=\frac{1}{n^{2}}\left(\mathbb{E}\left[\left(N_{d}\right)^{2}\right]-\mathbb{E}\left[\frac{1}{n}N_{d}\right]^{2}\right)$$

$$=\frac{1}{n^{2}}\left(\mathbb{E}\left[\left(N_{d}\right)^{2}\right]-\mathbb{E}\left[N_{d}\right]^{2}\right)$$

$$=\left[\left(\sum_{v\in V}\mathbb{I}\left[\operatorname{deg}(v)=d\right]\right)^{2}\right] (\operatorname{see Step 1})$$

$$N_{d}=\sum_{v\in V}\mathbb{I}_{\left\{\operatorname{deg}(v)=d\right\}} =\mathbb{E}\left[\left(\sum_{v\in V}\mathbb{I}\left\{\operatorname{deg}(v)=d\right\}\right)^{2}\right]$$

$$=\mathbb{E}\left[\sum_{v\in V}(\mathbb{1}\left\{\operatorname{deg}(v)=d\right\}\right)^{2}+\sum_{v\in V}\sum_{u\neq v}\mathbb{1}\left\{\operatorname{deg}(v)=d\right\}\cdot\mathbb{1}\left\{\operatorname{deg}(u)=d\right\}\right]$$

$$\operatorname{Indicator RV} X: X^{2}=X, =\mathbb{E}\left[\sum_{v\in V}\mathbb{I}\left\{\operatorname{deg}(v)=d\right\}\right]+\mathbb{E}\left[\sum_{v\in V}\sum_{u\neq v}\mathbb{1}\left\{\operatorname{deg}(v)=d\right\}\cdot\mathbb{1}\left\{\operatorname{deg}(u)=d\right\}\right]$$

$$\operatorname{Indicator RV} X: X^{2}=X, =\mathbb{E}\left[\sum_{v\in V}\mathbb{I}\left\{\operatorname{deg}(v)=d\right\}\right]+\mathbb{E}\left[\sum_{v\in V}\sum_{u\neq v}\mathbb{I}\left\{\operatorname{deg}(v)=d\right\}\cdot\mathbb{1}\left\{\operatorname{deg}(u)=d\right\}\right]$$

$$\operatorname{Ino of Exp.} =\sum_{v\in V}\mathbb{E}\left[\mathbb{1}\left\{\operatorname{deg}(v)=d\right\}\right]+\sum_{v\in V}\sum_{u\neq v}\mathbb{E}\left[\mathbb{1}\left\{\operatorname{deg}(v)=d\right\}\cdot\mathbb{1}\left\{\operatorname{deg}(u)=d\right\}\right]$$

$$=\operatorname{Pr}\left[\operatorname{deg}(v)=d\right]+\operatorname{Pr}\left[\operatorname{deg}(v)=d\right]+\operatorname{Pr}\left[\operatorname{deg}(v)=d\right]$$



$$\Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_{d}\right] - \frac{1}{n}N_{d}\right| \ge \varepsilon\right] \le \operatorname{Var}\left[\frac{1}{n}N_{d}\right]/\varepsilon^{2}$$

$$\operatorname{Var}\left[\frac{1}{n}N_{d}\right] = \mathbb{E}\left[\left(\frac{1}{n}N_{d}\right)^{2}\right] - \mathbb{E}\left[\frac{1}{n}N_{d}\right]^{2}$$

$$= \frac{1}{n^{2}}\left(\mathbb{E}\left[\left(N_{d}\right)^{2}\right] - \mathbb{E}\left[N_{d}\right]^{2}\right)$$

$$= \frac{1}{n^{2}}(n\operatorname{Pr}\left[\deg(v) = d\right]$$

$$+ n(n-1)\operatorname{Pr}\left[\deg(v) = d \wedge \deg(u) = d\right]$$

$$- (n\operatorname{Pr}\left[\deg(v) = d\right])^{2}\right)$$

$$= \frac{1}{n}\operatorname{Pr}\left[\deg(v) = d\right] \le 1$$

$$- \operatorname{Pr}\left[\deg(v) = d\right]^{2}$$

$$\le \frac{1}{n} + \operatorname{Pr}\left[\deg(v) = d \wedge \deg(u) = d\right]$$

$$- \operatorname{Pr}\left[\deg(v) = d\right]^{2}$$

$$\lim_{n\to\infty} \Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_d\right] - \frac{1}{n}N_d\right| \ge \varepsilon\right] = 0$$

**Chebychev**: X finite variance, b > 0 $\Pr[|X - \mathbb{E}[X]| \ge b] \le \operatorname{Var}[X]/b^2$ 

$$\left(\sum_{i}a_{i}\right)^{2}=\sum_{i}a_{i}^{2}+\sum_{i}\sum_{j\neq i}a_{i}a_{j}$$



$$\begin{aligned} & \Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_{d}\right]-\frac{1}{n}N_{d}\right|\geq\varepsilon\right] \leq \operatorname{Var}\left[\frac{1}{n}N_{d}\right]/\varepsilon^{2} \\ & \operatorname{Var}\left[\frac{1}{n}N_{d}\right]=\mathbb{E}\left[\left(\frac{1}{n}N_{d}\right)^{2}\right]-\mathbb{E}\left[\frac{1}{n}N_{d}\right]^{2} \\ & =\frac{1}{n^{2}}\left(\mathbb{E}\left[\left(N_{d}\right)^{2}\right]-\mathbb{E}\left[N_{d}\right]^{2}\right) \\ & \leq\frac{1}{n}+\Pr\left[\deg(v)=d\wedge\deg(u)=d\right] \\ & -\Pr\left[\deg(v)=d\right]\Pr\left[\deg(u)=d\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & -\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & -\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & \leq\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{2}+Y_{2}=d\right] \\ & \leq\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & \leq\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\right] \\ & \int \operatorname{Couplings} \\ & =\operatorname{Consider}\operatorname{deg}(u) \operatorname{and}\operatorname{deg}(v) \\ & =\operatorname{deg}(v) \operatorname{degendent} \operatorname{deg}(u) \\ & =\operatorname{deg}(v) \operatorname{degendent} \operatorname{deg}(v) \\ & =\operatorname{deg}(v)$$

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$$\begin{aligned} & \Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_{d}\right]-\frac{1}{n}N_{d}\right|\geq\varepsilon\right] \leq \operatorname{Var}\left[\frac{1}{n}N_{d}\right]/\varepsilon^{2} \\ & \operatorname{Var}\left[\frac{1}{n}N_{d}\right]=\mathbb{E}\left[\left(\frac{1}{n}N_{d}\right)^{2}\right]-\mathbb{E}\left[\frac{1}{n}N_{d}\right]^{2} \\ & =\frac{1}{n^{2}}\left(\mathbb{E}\left[\left(N_{d}\right)^{2}\right]-\mathbb{E}\left[N_{d}\right]^{2}\right) \\ & \leq\frac{1}{n}+\Pr\left[\deg(v)=d\wedge\deg(u)=d\right] \\ & -\Pr\left[\deg(v)=d\right]\Pr\left[\deg(u)=d\right] \\ & \deg(v)\stackrel{d}{=}\deg(u) \end{aligned} \end{aligned} \qquad \begin{aligned} & \left(\sum_{i}a_{i}\right)^{2}=\sum_{i}a_{i}^{2}+\sum_{i}\sum_{j\neq i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}=\sum_{i}a_{i}^{2}+\sum_{i}\sum_{i}\sum_{j\neq i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}\sum_{j\neq i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}=\sum_{i}a_{i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}\sum_{j\neq i}a_{j}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}\sum_{j\in i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}a_{i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}a_{i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}a_{i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{2}+\sum_{i}a_{i}a_{i}a_{i}a_{i}a_{j} \\ & \left(\sum_{i}a_{i}\right)^{$$

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$$\begin{aligned} & \Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_{d}\right]-\frac{1}{n}N_{d}\right|\geq\varepsilon\right] \leq \operatorname{Var}\left[\frac{1}{n}N_{d}\right]/\varepsilon^{2} \xrightarrow{n\to\infty} 0 \\ & \operatorname{Var}\left[\frac{1}{n}N_{d}\right]=\mathbb{E}\left[\left(\frac{1}{n}N_{d}\right)^{2}\right]-\mathbb{E}\left[\frac{1}{n}N_{d}\right]^{2}\leq\frac{1}{n}+2p=\frac{1}{n}+2\frac{c}{n}\xrightarrow{n\to\infty} 0 \\ & =\frac{1}{n^{2}}\left(\mathbb{E}\left[\left(N_{d}\right)^{2}\right]-\mathbb{E}\left[N_{d}\right]^{2}\right) \\ & \leq\frac{1}{n}+\Pr\left[\deg(v)=d\wedge\deg(u)=d\right] \\ & -\Pr\left[\deg(v)=d\right]\Pr\left[\deg(u)=d\right] \\ & -\Pr\left[\deg(v)=d\right]\Pr\left[\deg(u)=d\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\wedge X_{2}+Y_{2}\neq d\right] \\ & \leq\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\wedge X_{2}+Y_{2}\neq d\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}+Y_{1}=d\wedge X_{1}+Y_{2}=d\wedge X_{2}+Y_{2}\neq d|X_{1}=0\right]\Pr\left[X_{1}=0\right] \\ & =\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}+Y_{2}\neq d|X_{1}=0\right]\Pr\left[X_{1}=1\right] \\ & \leq\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}+Y_{2}\neq d|X_{1}=0\right] \\ & =\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}+Y_{2}\neq d|X_{1}=0\right] \\ & =\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}+Y_{2}\neq d|X_{1}=0\right] \\ & =\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}=1\right]\Pr\left[X_{1}=1\right] \\ & \leq\frac{1}{n}+\Pr\left[X_{2}=1\right]+\Pr\left[X_{1}=1\right] \\ & =\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}=1\right]\Pr\left[X_{1}=1\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}=1\right] \\ & =\frac{1}{n}+\Pr\left[Y_{1}=d\wedge Y_{2}=d\wedge X_{2}=1\right]\Pr\left[X_{1}=1\right] \\ & =\frac{1}{n}+\Pr\left[X_{1}=1\right] \\ & =\frac{1}{n}$$

## **Application: ER – Degree Distribution**

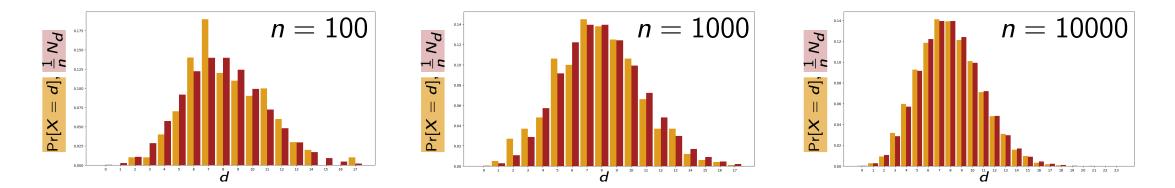
**Theorem:** Consider a G(n, p) with p = c/n for constant c > 0. For  $\lambda = -n \log(1-p)$ , let  $X \sim \text{Pois}(\lambda)$ . Then for all d > 0 and every  $\varepsilon > 0$  we have  $\lambda = c + O(1/n) \rightarrow c$  for  $n \rightarrow \infty$  $\lim_{n \rightarrow \infty} \Pr\left[\left|\Pr[X = d] - \frac{1}{n}N_d\right| \ge \varepsilon\right] = 0.$ 

#### Proof

Step 1: 
$$\Pr[X = d]$$
 is close to the expectation of  $\frac{1}{n}N_d$   $\lim_{n \to \infty} \left|\Pr[X = d] - \mathbb{E}\left[\frac{1}{n}N_d\right]\right| = 0$ 

• Step 2:  $\frac{1}{n}N_d$  is concentrated (via Chebychev)

$$\lim_{n\to\infty} \Pr\left[\left|\mathbb{E}\left[\frac{1}{n}N_d\right] - \frac{1}{n}N_d\right| \ge \varepsilon\right] = 0 \checkmark$$



## **Concentration Bounds So Far**



**Definition**: A **concentration inequality** bounds the probability of a random variable to deviate from a given value (typically its expectation) by a certain amount.

### Markov

- based on expectation (first moment)
- X non-negative random variable and a > 0 $\Pr[X > a] < \mathbb{E}[X]/a$

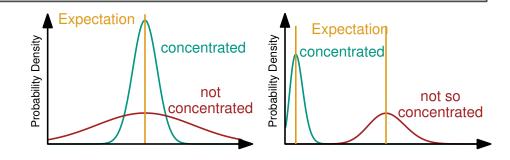
$$\Pr[X \ge a] \le$$

tight

## Chebychev

- based on variance (second moment)
- X random variable with finite variance and b > 0 $\Pr[|X - \mathbb{E}[X]| \ge b] \le \operatorname{Var}[X]/b^2$
- tight (stated without proof)

Can we utilize higher-order moments for even stronger bounds?



## **Another Moment Please**



The *n*-th raw moment of a random variable X is  $\mathbb{E}[X^n]$ 

• We can capture *all* moments of *X* using a single function

Looks scary, but is again just  $\mathbb{E}[f(X)]$  for  $f(X) = e^{tX}$ 

**Markov**: X non-negative, b > 0:

 $|\Pr[X \ge b] \le \mathbb{E}[X]/b.$ 

**Definition**: For a random variable X the moment generating function is  $M_X(t) = \mathbb{E}[e^{tX}]$ 

• Where the name comes from: For the *n*-th derivative  $M_X^{(n)}(t)$  we have  $M_X^{(n)}(0) = \mathbb{E}[X^n]$ 

**Theorem**: For independent random variables  $X, Y: M_{X+Y}(t) = M_X(t) \cdot M_Y(t)$ .

**Proof** 
$$M_{X+Y}(t) = \mathbb{E}[e^{t(X+Y)}] = \mathbb{E}[e^{tX} \cdot e^{tY}] = \mathbb{E}[e^{tX}] \cdot \mathbb{E}[e^{tY}] = M_X(t) \cdot M_Y(t)$$

**Concentration Inequality** Had his 100th birthday in 2023! Thought the bound (now named after him) to be so trivial that he didn't mention that it actually came from Herman Rubin. (A conversation with Herman Chernoff, John Bather, Statist. Sci. 1996)

**Theorem (Chernoff Bounds)**: Let X be a random variable and a > 0. Then,  $\Pr[X \ge a] \le \min_{t>0} \mathbb{E}[e^{tX}]/e^{ta}$  and  $\Pr[X \le a] \le \min_{t<0} \mathbb{E}[e^{tX}]/e^{ta}$ .

**Proof** for all t > 0:  $\Pr[X \ge a] = \Pr[e^{tX} \ge e^{ta}] \le \mathbb{E}[e^{tX}]/e^{ta}$  $\checkmark < \min_{t>0} \mathbb{E}[e^{tX}]/e^{ta} \checkmark$ 

for all t < 0: analogous.  $\checkmark$ 

Get bounds for specific random variables by finding a good t!



## **Application: Binomial Distribution**

$$\begin{array}{l} \textbf{Theorem: Let } X \sim \text{Bin}(n, p). \text{ Then for any } \varepsilon > 0 \\ & \Pr[X \ge (1+\varepsilon)\mathbb{E}[X]] \le \left(\frac{e^{\varepsilon}}{(1+\varepsilon)^{(1+\varepsilon)}}\right)^{\mathbb{E}[X]} \\ \end{array} \\ \begin{array}{l} \textbf{Proof Consider } X \text{ as the sum of independent } X_i \sim \text{Ber}(p) \\ & M_{X_i}(t) = \mathbb{E}[e^{tX_i}] = \Pr[X_i = 0] \cdot e^{t \cdot 0} + \Pr[X_i = 1] \cdot e^{t \cdot 1} \\ & = (1-p) + pe^t = 1 + (e^t - 1)p \le e^{(e^t - 1)p} \\ & 1+x \le e^{\times} \end{array} \\ \begin{array}{l} \textbf{Max}(t) = M_{\sum X_i}(t) = \prod_{i=1}^n M_{X_i}(t) \le \prod_{i=1}^n e^{(e^t - 1)p} \\ & e^{(e^t - 1)\mathbb{E}[X]} \end{array} \\ \textbf{Pr}[X \ge (1+\varepsilon)\mathbb{E}[X]] \le \min_{t>0} \frac{\mathbb{E}[e^{tX}]}{e^{t(1+\varepsilon)\mathbb{E}[X]}} \le \min_{t>0} \frac{e^{(e^t - 1)\cdot p}}{e^{t(1+\varepsilon)\mathbb{E}[X]}} = \min_{t>0} \left(\frac{e^{(e^t - 1)}}{e^{t(1+\varepsilon)}}\right)^{\mathbb{E}[X]} \\ \textbf{Fr}[X \ge (1+\varepsilon)\mathbb{E}[X]] \le \min_{t>0} \frac{\mathbb{E}[e^{tX}]}{e^{t(1+\varepsilon)\mathbb{E}[X]}} \le \min_{t>0} \frac{e^{(e^t - 1)\mathbb{E}[X]}}{e^{t(1+\varepsilon)\mathbb{E}[X]}} = \min_{t>0} \left(\frac{e^{(e^t - 1)}}{e^{t(1+\varepsilon)}}\right)^{\mathbb{E}[X]} \\ \textbf{Fr}[X \ge 16] = \Pr[X \ge (1+3)\mathbb{E}[X]] \le \left(\frac{e^3}{(1+3)^{1+3}}\right)^4 = \frac{e^{1^2}}{4^{4^4}} \approx 0.00003789 \end{aligned}$$



## **Chernoff – Simpler Versions**

**Theorem:** Let 
$$X \sim Bin(n, p)$$
. Then for any  $\varepsilon > 0$   
 $Pr[X \ge (1 + \varepsilon)\mathbb{E}[X]] \le \left(\frac{e^{\varepsilon}}{(1 + \varepsilon)^{(1 + \varepsilon)}}\right)^{\mathbb{E}[X]}$ .  
  
**Chernoff:** Random variable X and  $a > 0$ :  
 $Pr[X \ge a] \le \min_{t>0} \mathbb{E}[e^{tX}]/e^{ta}$ .

**Corollary**: Let  $X \sim Bin(n, p)$ . Then for any  $t \ge 6\mathbb{E}[X]$ ,  $Pr[X \ge t] \le 2^{-t}$ .

**Corollary**: Let  $X \sim Bin(n, p)$ . Then for any  $\varepsilon \in (0, 1]$ ,  $Pr[X \ge (1 + \varepsilon)\mathbb{E}[X]] \le e^{-\varepsilon^2/3 \cdot \mathbb{E}[X]}$ .

**Corollary**: Let  $X \sim Bin(n, p)$ . Then for any  $\varepsilon \in (0, 1)$ ,  $Pr[X \leq (1 - \varepsilon)\mathbb{E}[X]] \leq e^{-\varepsilon^2/2 \cdot \mathbb{E}[X]}$ .

In fact, these also work when the X<sub>i</sub> are Bernoulli random variables with different success probabilities

## Conclusion

#### Concentration

- Is a random variable likely to yield values close to the expectation?
- Concentration inequalities bound the probability for a random variable to deviate from its expectation

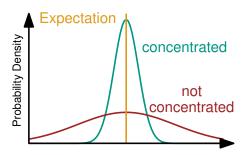
### Moments

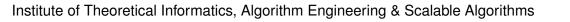
- Used to characterize the shape of a distribution
- First moment: expected value
- Second moment: variance
- Moment generating functions to determine higher-order moments

## **Concentration Inequalities**

- Markov: Based on first moment
- Chebychev: Squaring within Markov (utilizing second moment)
- Chernoff: Exponentiating within Markov (utilizing moment generating functions)
- Examples: Sum of coin flips, empirical degree distribution of ER graphs







 $\frac{1}{n}N_d$ 



n = 10000